

# The **DOT** Calculus

(**D**ependent **O**bject **T**ypes)

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# DOT: Dependent Object Types

- ▶ DOT is a core calculus for path-dependent types.
- ▶ Goals
  - ▶ simplify Scala's type system by desugaring to DOT
  - ▶ simplify Scala's type inference by relying on DOT
  - ▶ prove that DOT is type-safe

# Types in Scala and DOT

# Types in Scala

modular            named type `scala.collection.BitSet`

                  compound type `Channel with Logged`

                  refined type `Channel { def close(): Unit }`

functional    parameterized type `List[String]`

                  existential type `List[T] forSome { type T }`

                  higher-kinded type `List`

## Reducing Functional to Modular?

- ▶ type parameter to type member

```
class List[Elem] {} /*vs*/ class List { type Elem }
```

- ▶ parameterized type to refined type

```
List[String] /*vs*/ List { type Elem = String }
```

- ▶ existential type?

```
List[_] /*or*/ List[T] forSome { type T } /*~vs~*/ List
```

- ▶ higher-kinded type?

```
List /*~vs~*/ List
```

## Bonus: Don't Repeat Yourself

```
        new HashMap[String, List[Int]]  
with SynchronizedMap[String, List[Int]]
```

```
    /*vs*/
```

```
        new HashMap[String, List[Int]]  
with SynchronizedMap
```

## Type Members and Modularity (Example)

```
trait Symbols {  
  type Type  
  // a symbol has a type ...  
  trait Symbol { def tpe: Type }  
}  
trait Types {  
  type Symbol  
  // ... and a type has a symbol  
  trait Type { def sym: Symbol }  
}  
// ... but the references are not hard-coded!  
// putting the cake together:  
trait SymbolTable extends Symbols with Types
```

## Type Members and Data Abstraction (Example)

```
trait HeapModule {  
  type Heap  
  type Elem  
  def ord: Ordering[Elem]  
  def empty: Heap  
  def isEmpty(h: Heap): Boolean  
  def insert(x: Elem, h: Heap): Heap  
  def findMin(h: Heap): Elem  
  def deleteMin(h: Heap): Heap  
  def merge(h1: Heap, h2: Heap): Heap  
}  
  
trait BinomialHeapModule extends HeapModule {  
  type Rank = Int  
  case class Node(x: Elem, r: Rank, c: List[Node])  
  override type Heap = List[Node]  
  // ... implementation ...  
}
```



## Path-Dependent Types and Binary Methods (Example)

```
val m1: HeapModule = ...
val m2: HeapModule = ...
// in Scala REPL
> m1.merge(m1.empty, m2.empty)
// error: type mismatch;
// found   : m2.Heap
// required: m1.Heap
//           m1.merge(m1.empty, m2.empty)
//                                     ^
```

# Types in DOT

types S, T, U

path-dependent type p.L

refined type T { z => D }

intersection T & T

union T | T

top Any

bottom Nothing

declarations D

type declaration type L >: S <: U

field declaration val l: U

method declaration def m(x: S): U

# Type Inference

## Type Inference in Scala (Least Upper Bound?)

```
trait A { type T <: A }
trait B { type T <: B }
trait C extends A with B { type T <: C }
trait D extends A with B { type T <: D }
// in Scala, lub(C, D) is an infinite sequence
A with B { type T <: A with B { type T <: ... } }

// in Scala REPL
> val o = if (true) (new C{}) else (new D{})
o: A with B{type T <: A with B} = ...
> val o:A with B{type T<:A with B {type T<:A with B}} =
    if (true) (new C{}) else (new D{})
o: A with B{type T <: A with B{type T <: A with B}} = ...
```





## Type Inference in Scala (Working too Hard too Soon?)

```
import scala.collection.mutable.{Map => MMap, Set => MSet}  
val ms: MMap[Int, MSet[Int]] = MMap.empty
```

```
// in Scala REPL
```

```
> if (!ms.contains(1)) ms += 1 -> MSet(1) else ms(1) += 1
```

```
res0: ... (796 characters) ...
```

```
> :t res0
```

```
: ... (21481 characters) ...
```

- ▶ Inspired by a bug report  
SI-5862: very slow compilation due to humonguous LUB.
- ▶ The character lengths reported are for Scala 2.11 (*after* the fix).
- ▶ In Dotty, type inference can be lazy thanks to the native unions (for least upper bounds) and intersections (for greatest lower bounds) of the core calculus (DOT).

# Meta-Theory



## Type Safety: “Well-typed programs can’t go wrong.”

```
// ... at least in some ways ...  
// JavaScript REPL  
> var x = {foo: 'hello'}  
undefined  
> x.foo  
"hello"  
> x.bar  
undefined  
> x.bar.boo  
// TypeError: Cannot read property 'boo' of undefined  
> x.foo('hello')  
// TypeError: string is not a function
```

# Type Safety: Syntactic Recipe (Concepts)

**small-step** operational semantics (“show all your steps”)

**irreducible** terms, stuck terms vs values

**progress** theorem: if a term type-checks, then it can take a step or it is a value (but not stuck!)

**preservation** theorem: if a term type-checks and steps, then the new term also type-checks with the same type (or subtype)

**type safety** = progress + preservation

## Preservation Challenge: Branding

```
trait Brand {
  type Name
  def id(x: Any): Name
}
// in REPL
val brand: Brand = new Brand {
  type Name = Any
  def id(x: Any): Name = x
}
brand.id("hi"): brand.Name // ok
"hi": brand.Name // error but probably sound
val brandi: Brand = new Brand {
  type Name = Int
  def id(x: Any): Name = 0
}
brandi.id("hi"): brandi.Name // ok
"hi": brandi.Name // error and probably unsound
```

## Theory: Subtyping of Path-Dependent Types

- ▶ `/*If*/ p.L /*has lower bound*/ Sp /*and upper bound*/ Up`
  - ▶ `S <: p.L /*if*/ S <: Sp`
  - ▶ `p.L <: U /*if*/ Up <: U`

## Preservation Challenge: Path equality

```
trait B { type X; val l: X }
val b1: B = new B { type X = String; val l: X = "hi" }
val b2: B = new B { type X = Int;    val l: X = 0    }
trait A { val i: B }
val a = new A { val i: B = b1 }
println(a.i.l : a.i.X) // ok
println(b1.l  : b1.X)  // ok
println(b1.l  : a.i.X) // error: type mismatch;
// found    : b1.l.type (with underlying type b1.X)
// required: a.i.X
// abstractly, would need to show
Any <: Nothing // lattice collapse!
// to show
b1.X <: a.i.X
// because
U of b1.X = Any <: Nothing = S of a.i.X
println(b2.l  : a.i.X) // error and probably unsound
```

## Challenge: Subtyping Transitivity

► `S <: T <: U /*imply*/ S <: U`

## Subtyping Transitivity and Path-Dependent Types

- ▶  $S <: p.L <: U$
- ▶  $p.L$  /\*has lower bound\*/  $S_p$  /\*and upper bound\*/  $U_p$
- ▶  $S <: S_p$
- ▶  $U_p <: U$

# Subtyping Transitivity and Path-Dependent Types with Bad Bounds

- ▶ `/*Suppose*/ p.L /*has bad bounds, e.g.*/`
- ▶ `p.L /*has lower bound*/ Any /*and upper bound*/ Nothing`
- ▶ `S <: Any /*for any*/ S`
- ▶ `U <: Nothing /*for any*/ U`
- ▶ `S <: p.L <: U /*imply*/ S <: U /*for any*/ S, U`
- ▶ Lattice collapse!



## Path-Dependent Types (Recap Example)

```
trait Animal { type Food; def gets: Food
                def eats(food: Food) {}; }
trait Grass; trait Meat
trait Cow extends Animal with Meat {
  type Food = Grass; def gets = new Grass {} }
trait Lion extends Animal {
  type Food = Meat; def gets = new Meat {} }
val leo = new Lion {}
val milka = new Cow {}
leo.eats(milka) // ok
val lambda: Animal = milka
lambda.eats(milka) // type mismatch
// found : Cow
// required: lambda.Food
lambda.eats(lambda.gets) // ok
```

## Path-Dependent Types (Recap Example Continued)

```
def share(a1: Animal)(a2: Animal)
  (bite: a1.Food with a2.Food) {
  a1.eats(bite)
  a2.eats(bite)
}

val simba = new Lion {}
share(leo)(simba)(leo.gets) // ok
share(leo)(lambda)(leo.gets) // error: type mismatch
// found : Meat
// required: leo.Food with lambda.Food

// Observation:
// We don't know whether the parameter type of
share(lambda1)(lambda2)_
// is realizable until run-time.
```

## Realizability of Intersection Types at Run-Time

```
val lambda1: Animal = new Lion {}
val lambda2: Animal = new Cow  {}
lazy val p: lambda1.Food & lambda2.Food = ???
// for illustration, say we re-defined the following:
trait Food { type T }
trait Meat extends Food { type T = Nothing }
trait Grass extends Food { type T = Any }
// statically
p.T /*has fully abstract bounds*/
// at runtime
lambda1.Food /*is*/ Meat /*&*/ lambda2.Food /*is*/ Grass
p /*has type*/ Meat & Grass
// lower bound is union of lower bounds
p.T /*has lower bound*/ Nothing | Any /*is*/ Any
// upper bound is intersection of upper bounds
p.T /*has upper bound*/ Nothing & Any /*is*/ Nothing
p.T /*has bad bounds at run-time!*/
```

# Summary

- ▶ DOT is a core calculus for path-dependent types.
- ▶ Benefits for Scala
  - ▶ simpler type system
  - ▶ better type inference
  - ▶ (hopefully) easier reasoning
- ▶ Challenges in Meta-Theory
  - ▶ type preservation
    - ▶ run-time path equality
    - ▶ inlining branding
  - ▶ path-dependent types with “bad bounds”
    - ▶ subtyping transitivity
    - ▶ run-time “bad bounds” via intersection types

Thank You!